Graph-Based Resource Allocation with Conflict Avoidance for V2V Broadcast Communications

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 3GPP¹ recently proposed a novel resource allocation notion called vehicle-to-vehicle (V2V) mode-3.

¹The 3rd Generation Partnership Project

²Initially aimed at supporting proximity services (ProSe).

³Pilot symbols more closely spaced for channel estimation in high Doppler



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- 3GPP¹ recently proposed a novel resource allocation notion called vehicle-to-vehicle (V2V) mode-3.
- V2V is based on a previously developed technology, namely device-to-device (D2D) communications²

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- 3GPP¹ recently proposed a novel resource allocation notion called vehicle-to-vehicle (V2V) mode-3.
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- Additional modifications have been applied in order to support more dynamic scenarios
 - Denser distribution of DMRS³

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- V2V is based on a previously developed technology, namely device-to-device (D2D) communications²
- Additional modifications have been applied in order to support more dynamic scenarios
 - Denser distribution of DMRS³
 - A novel structure that supports adjacent (i) scheduling assignments and (ii) data resources

³Pilot symbols more closely spaced for channel estimation in high Doppler



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 Thus, besides uplink and downlink (Uu), vehicles can also communicate via sidelink (PC5), which sustains direct communications between vehicles.



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 Conversely to mainstream communications, in V2V mode-3 data traffic from/to vehicles do not traverse the eNodeB.



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 - Thereupon, vehicles communicate directly—with their counterparts via sidelink—in a broadcast manner.

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- Thus, in V2V *mode-3* operation⁴:
 - eNodeBs only intervene in the resource allocation process.
 - Thereupon, vehicles communicate directly—with their counterparts via sidelink—in a broadcast manner.
- In safety applications, vehicles would typically exchange cooperative awareness messages (CAMs).

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A CAM message contains information such as

- position,
- velocity,
- direction, etc.
- of a vehicle.



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• As these messages transport important information, it is crucial that they are transmitted **reliably**.



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- Due to the one-to-all nature of V2V mode-3, the allocation of resources (or subchannels) slightly differs from mainstream communications.

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- As these messages transport important information, it is crucial that they are transmitted **reliably**.
- Due to the one-to-all nature of V2V mode-3, the allocation of resources (or subchannels) slightly differs from mainstream communications.
- **Example:** If two vehicles transmit concurrently they will not receive the CAM message of the other.



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Vehicular Scenario



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Vehicular Scenario



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Sidelink Channelization



- T: duration of a subframe
- K: number of subchannels per subframe
- L: total number of subframes for allocation
- B: subchannel bandwidth



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Subchannel Structure

Assuming a 10 MHz ITS (Intelligent Transportation Systems) channel, up to 7 subchannels per subframe can be obtained. Thus,

- B: 1.26 MHz
- *T*: 1 ms (2 slots of 0.5 ms each)
- Control: 2 RBs⁵ per slot \leftarrow 24 subcarriers
- Data: 5 RBs per slot \leftarrow 60 subcarriers

Subchannel

A subchannel of 7 RBs is capable of transporting a basic CAM message with a payload of 200 bytes.

⁵RB: A resource block consits of 12 subcarriers

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Problem Formulation

Antecedents

- Vehicles can either transmit or receive at a time⁶.
- When two or more vehicles transmit concurrently in subchannels of the same subframe, a conflict is generated.

Objectives

- Attain a conflict-free subchannel assignment.
- Maximize the sum-capacity of the system



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Problem Formulation (cont'd)

Antecedents

- Vehicles can either transmit or receive at a time.
- When two or more vehicles transmit concurrently in subchannels of the same subframe, a conflict is generated.

Proposed Solution

- The subchannel allocation problem is approached as a bipartite graph matching.
- Additional constraints have been considered in order to prevent conflicts from occurring.



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There are available spectrum resources for uplink and downlink.



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- There are available spectrum resources for uplink and downlink.
- Vehicles report to eNodeBs the channel conditions they perceive (e.g. CQI, SINR).



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- Vehicles report to eNodeBs the channel conditions they perceive (e.g. CQI, SINR).
- The eNodeB performs the assignment of subchannels based on the information received.
- The eNodeB notifies the vehicles on their assigned subchannel via downlink.
- Then, vehicles start broadcasting CAM messages.



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Graph Representation of Subchannel Allocation

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The optimization problem can be recast as:

max
$$\mathbf{c}^T \mathbf{x}$$
 (2a)
subject to $\left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \otimes \mathbf{1}_{1 \times K} \mathbf{x} = \mathbf{1}$ (2b)

<u>Note:</u> For completeness, we have assumed that the number of vehicles is equal to the number of subframes, i.e. N = L

 r_j This problem structure cannot be exploited to be approached by known matching algorithms. So we proceed as follows

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(1+SINRij)

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Objective Function

 $\max \mathbf{c}^T \mathbf{x}$

Because $\mathbf{x} \in \mathbb{B}^{MK}$, then the objective function can be recast as

$$\mathbf{c}^T \mathbf{x} \equiv \mathbf{x}^T diag(\mathbf{c}) \mathbf{x}$$

without affecting optimality. Note that $M = N^2$.



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Objective Function

 $\max \mathbf{c}^T \mathbf{x}$

For any vehicle v_i ,

$$x_{ij}x_{ik} = 0, \ r_j, r_k \in \mathcal{R}_{\alpha}.$$

Moreover,

$$c_{ij}x_{ij}x_{ik} = 0, \ r_j, r_k \in \mathcal{R}_{\alpha}.$$

In general, for N vehicles

$$\mathbf{x}^{T} \big(\mathbf{I}_{M \times M} \otimes [\mathbf{1}_{K \times K} - \mathbf{I}_{K \times K}] \big) diag(\mathbf{c}) \mathbf{x} = 0.$$

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 $\max \mathbf{c}^T \mathbf{x}$

As long as $\mathbf{x}^T (\mathbf{I}_{M \times M} \otimes [\mathbf{1}_{K \times K} - \mathbf{I}_{K \times K}]) diag(\mathbf{c}) \mathbf{x} = 0$ holds, conflicts will be prevented.

We can aggregate this condition to the objective function. Hence,

$$\mathbf{c}^T \mathbf{x} = \mathbf{x}^T diag(\mathbf{c}) \mathbf{x} + \mathbf{x}^T \big(\mathbf{I}_{M \times M} \otimes [\mathbf{1}_{K \times K} - \mathbf{I}_{K \times K}] \big) diag(\mathbf{c}) \mathbf{x}$$

Further manipulation leads to

$$\mathbf{c}^T \mathbf{x} = \mathbf{x}^T (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{K \times K}) diag(\mathbf{c}) \mathbf{x}$$

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 $\max \mathbf{c}^T \mathbf{x}$

Property 1 (Product of two tensor products) Let $\mathbf{X} \in \mathbb{R}^{m \times n}$, $\mathbf{Y} \in \mathbb{R}^{r \times s}$, $\mathbf{W} \in \mathbb{R}^{n \times p}$, and $\mathbf{Z} \in \mathbb{R}^{s \times t}$, then $\mathbf{X}\mathbf{Y} \otimes \mathbf{W}\mathbf{Z} = (\mathbf{X} \otimes \mathbf{W})(\mathbf{Y} \otimes \mathbf{Z}) \in \mathbb{R}^{mr \times pt}$ $\mathbf{c}^T \mathbf{x} = \mathbf{x}^T (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{K \times K}) diag(\mathbf{c}) \mathbf{x}$ $= \mathbf{x}^{T} (\mathbf{I}_{M \times M} \mathbf{I}_{M \times M} \otimes \mathbf{1}_{K \times 1} \mathbf{1}_{1 \times K}) diag(\mathbf{c}) \mathbf{x}$ $= \mathbf{x}^T (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{K \times 1}) (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{1 \times K}) diag(\mathbf{c}) \mathbf{x}$ TU/e \mathbf{y}^T

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Constraints

subject to
$$\left[\begin{array}{c} \mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N} \\ \hline \mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N} \end{array} \right] \otimes \mathbf{1}_{1 \times K} \mathbf{x} = \mathbf{1}$$

Property 2 (Pseudo-inverse of a tensor product) Let $\mathbf{X} \in \mathbb{R}^{m \times n}$ and $\mathbf{Y} \in \mathbb{R}^{r \times s}$, then $(\mathbf{X} \otimes \mathbf{Y})^{\dagger} = \mathbf{X}^{\dagger} \otimes \mathbf{Y}^{\dagger} \in \mathbb{R}^{ns \times mr}$

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Constraints

subject to
$$\left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \otimes \mathbf{1}_{1 \times K} \mathbf{x} = \mathbf{1}$$

$$\begin{pmatrix} \left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \otimes \mathbf{1}_{1 \times K} \end{pmatrix} \begin{pmatrix} \mathbf{I}_{M \times M} \otimes \mathbf{1}_{1 \times K}^{\dagger} \end{pmatrix} \mathbf{y} = \mathbf{1}$$

$$= \begin{pmatrix} \left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \mathbf{I}_{M \times M} \end{pmatrix} \otimes \underbrace{\begin{pmatrix} \mathbf{1}_{1 \times K} \mathbf{1}_{1 \times K}^{\dagger} \end{pmatrix}}_{\mathbf{1}} \mathbf{y} = \mathbf{1}$$

$$= \begin{bmatrix} \frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \end{bmatrix} \mathbf{y} = \mathbf{1}$$

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Original Problem

$$\max \mathbf{c}^T \mathbf{x}, \text{ subject to } \left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \otimes \mathbf{1}_{1 \times K} \mathbf{x} = \mathbf{1}$$

Resultant Problem

$$\max \mathbf{d}^T \mathbf{y}, \qquad \text{subject to } \left[\frac{\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times N}}{\mathbf{1}_{1 \times N} \otimes \mathbf{I}_{N \times N}} \right] \mathbf{y} = \mathbf{1}.$$

where $\mathbf{d} = (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{1 \times K}) diag(\mathbf{c}) \mathbf{x}$ and $\mathbf{y} = (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{1 \times K}) \mathbf{x}$ Dimensionality reduction: $\rightarrow |\mathbf{x}| = MK \quad \rightarrow |\mathbf{y}| = M$. The resultant problem can now be approached through the

Kuhn-Munkres Algorithm.

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Transformation



$$\mathbf{d} = \lim_{\beta \to \infty} \frac{1}{\beta} \log^{\circ} \left\{ (\mathbf{I}_{M \times M} \otimes \mathbf{1}_{1 \times K}) \mathrm{e}^{\circ \beta \mathbf{c}} \right\}$$

 $\log^{\circ}{\{\cdot\}}$: Element-wise natural logarithm. $e^{\circ\{\cdot\}}$ Hadamard exponential.

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Simulations: Data Rate per Vehicle

Rate [Mbps / subchannel] 15Exhaustive Search Graph-based Algorithm Greedy Algoritm Random Algorithm 8.97 8.97 8.97 8.22 8.22 8.02 7.63 107.12 7.12 85 4.525.76 1.67.25 .16.160 **Highest-Rate** Worst-Rate System System Rate Vehicle Vehicle Average Rate Std. Dev. Description Value Number of vehicles per cluster 100 Number of clusters 4 Message rate (Hz) 10

Number of allottable subframes 100 Number of resources per subframe 7

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Simulations: Least Favored Vehicle

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Simulations: CDF of Rate Values

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Simulations: One-shot simulation

Rate [Mbps / subchannel] 15Vehicle v_4 Vehicle v1 Vehicle v_2 Vehicle v_3 Vehicle v_5 Vehicle v_6 Vehicle v7 Vehicle v₈ Vehicle vo Vehicle v_{10} 04 1088 r-66 50 Exhaustive Greedy Graph-based Random Search Algorithm Algorithm Algorithm Description Value Number of vehicles per cluster 10 Number of clusters 1 Message rate (Hz) 10 Number of allottable subframes 10 3 Number of resources per subframe

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Complexity

- Exhaustive search: $\mathcal{O}(|\mathcal{R}|!/(|\mathcal{R}| |\mathcal{V}|)!)$
- Graph-based approach: $\mathcal{O}(\max\{|\mathcal{V}|, |\mathcal{R}|/K\}^3)$
- Greedy algorithm: $\mathcal{O}(|\mathcal{V}||\mathcal{R}|)$
- Random algorithm: $\mathcal{O}(|\mathcal{V}|)$
- $|\mathcal{V}|$: Number of vehicles
- $|\mathcal{R}|$: Number of resources
- K: Number of subchannels per subframe



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Conclusions



- A graph theoretical approach for subchannel allocation in V2V mode-3 was presented.
- Subchannel conflict avoidance was enforced through graph vertex aggregation.
- For the case of independent vehicular clusters, the proposed approach attains the same optimality as exhaustive search at lower complexity.
- Although not explicitly enforced, the proposed scheme is capable of improving the rate fairness between vehicles.

Questions



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